

1/3 Sun Aug 26 2007

~~Technically, any non-capturing~~

~~But for recursion, any non-capturing~~

Any non-capturing clause can be rewritten as a capturing clause + α -conversion.

$$[\{P \rightarrow V\}] = \text{let } t_1 = \text{self}, t_2 = \text{super} \\ \text{in } [P \rightarrow [V] \sigma_{\text{offspring}}]$$

\therefore ignore non-capturing clauses in the core calculus. Also ignore tuples + other complexity of pattern-matching

core λ : $M, N = \lambda x. M \mid M N \mid x$

core TNG: ~~$M, N = [P \rightarrow M] \mid M N \mid M \oplus N \mid x$~~
 ~~$P = \text{atom} \mid x$~~

$$M, N = [P \rightarrow M] \mid M N \mid M \oplus N \mid V \mid x$$

$$P = V \mid x$$

$$V = \text{atom}$$

~~Evaluation context of~~

~~exprs $M, N = [P \rightarrow M] \mid V \oplus M$~~

$E \sigma : M \rightarrow V$ *She.*

exprs $M, N = [P \rightarrow M] \mid M N \mid M \oplus N \mid L \mid x$
 pat1 $P = L \mid x$
 lib $L = \text{atom}$

values V, W ~~$[P \rightarrow M] \mid L \mid V \oplus W \mid \emptyset$~~

2/3

$$\forall v. v \equiv v \oplus \emptyset$$

$$\mathcal{E}\Gamma : M \rightarrow V, \quad \text{eval.}$$

$$\mathcal{E}\Gamma [P \rightarrow M] = [\Gamma \vdash P \rightarrow M]$$

$$\mathcal{E}\Gamma \ M \ N = \text{send}(\mathcal{E}\Gamma M, \mathcal{E}\Gamma N)$$

$$\mathcal{E}\Gamma \ M \oplus N = \mathcal{E}\Gamma M \oplus \mathcal{E}\Gamma N$$

$$\mathcal{E}\Gamma \ L = L$$

$$\mathcal{E}\Gamma \ x = \begin{cases} V & \text{if } x \vdash V \text{ in } \Gamma \\ \perp & \text{otherwise} \end{cases}$$

$$\Gamma = \emptyset \mid x \vdash V, \Gamma$$

$$\text{send}(L, V) = L \oplus V$$

$$\text{send}(L \oplus V, W) = \perp$$

Forgot self & super!

$$M, N = [P \rightarrow M] \mid M \ N \mid M \oplus N \mid L \mid x \mid \text{self} \mid \text{super}$$

$$P = L \mid x$$

$$L = \text{atom.}$$

$$V, W = [\Gamma \vdash P \rightarrow M] \mid L \mid \emptyset \mid V \oplus W$$

$$\Gamma = \emptyset \mid x \vdash V, \Gamma$$

$$\forall v: V. v \oplus \emptyset \equiv v$$

$$\mathcal{E}\Gamma_{s_1, s_2} [P \rightarrow M] = \mathcal{E}\Gamma_{s_1} [\Gamma \vdash P \rightarrow M]$$

$$M \ N = \text{send}(e M, e N)$$

$$M \oplus N = e M \oplus e N$$

$$L = L$$

$$x = \begin{cases} \Gamma x & \text{otherwise } \perp \end{cases}$$

$$\text{self} = s_1$$

$$\text{super} = s_2$$

$$\text{send} : (V, V) \rightarrow V$$

$$\text{send} (\emptyset \oplus V, -) \rightarrow \perp$$

$$\text{send} (L \oplus V, -) \rightarrow \perp$$

~~$$\text{send} ([\Gamma, s_1, s_2 \vdash P \rightarrow M] \oplus V, L) = \mathcal{E} \Gamma_{s_1, s_2} M$$~~

~~$$\text{send} ([\Gamma, s_1, s_2 \vdash L \rightarrow M] \oplus V, L) = \mathcal{E} \Gamma_{s_1, s_2} M$$~~

~~$$\text{send} ([\Gamma, s_1, s_2 \vdash$$~~

$$\text{send} (V_R, V_M) = \text{search } V_R \ V_M$$

$$\text{search } V_R \ \emptyset \oplus \emptyset \ V_M = \perp$$

$$\text{search } V_R \ L \oplus V_S \ V_M = \text{search } V_R \ V_S \ V_M$$

$$\text{search } V_R \ [\Gamma \vdash L \rightarrow M] \oplus V_S \ L = \mathcal{E} \Gamma \ V_R \ V_S \ M$$

$$\text{search } V_R \ [\Gamma \vdash x \rightarrow M] \oplus V_S \ V_M = \mathcal{E} (x \mapsto V_M, \Gamma) \ V_R \ V_S \ M$$

$$\text{search } V_R \ - \oplus V_S \ V_M = \text{search } V_R \ V_S \ V_M$$

Now, tuples.

$$\text{search } V_R \ [\Gamma \vdash P \rightarrow M] \ V_M = \mathcal{E} \Gamma_2 \ V_R \ V_S \ M$$

$$\text{where } \Gamma_2 = \Gamma_1 \uparrow \Gamma$$

$$\text{if } \text{match}(P, V_M) = \Gamma_1$$

$$\emptyset \uparrow \Gamma = \Gamma \quad ; \quad x \mapsto V, \Gamma_1 \uparrow \Gamma_2 = x \mapsto V, (\Gamma_1 \uparrow \Gamma_2)$$

$$\text{match}(L, L) = \emptyset$$

$$\text{match}(x, V) = x \mapsto V, \emptyset$$

$$\text{match}((P_1, P_2), (V_1, V_2)) = \Gamma_1 \uparrow \Gamma_2 \text{ if } \text{match}(P_1, V_1) = \Gamma_1 \\ \text{and } \text{match}(P_2, V_2) = \Gamma_2$$

Without tuples, not powerful enough to make simultaneous bindings, \therefore capturing self & super simultaneously not possible

— WRONG. Carrying + renaming works ok:

$$[\text{msg} : [\text{xself} = [\text{xsuper} = \dots]] \text{ self super msg}]$$

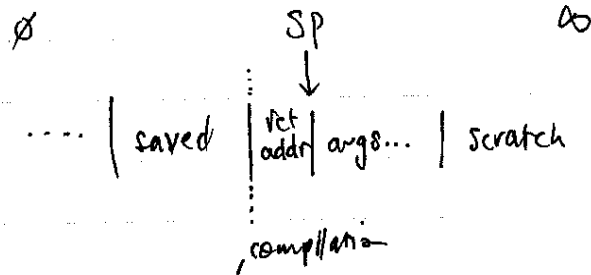
exp := [p → exp; ...]
 | exp exp
 | exp + exp
 | .atom
 | var

CLOSE
 SEND
 EXTEND
 LIT
 REF

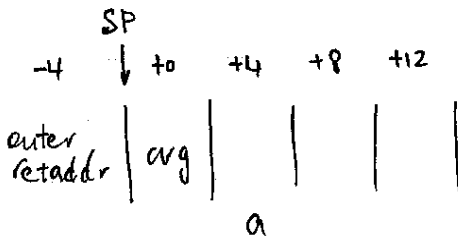
return
regs?
✓
callee
save?

ax accumulator
 ✓ bx
 cx argc?
 ✓ dx
 ✓ si
 ✓ di
 ✓ bp
 n/a sp

(a b) c
 a (b c)



[(a b) c] ⇒ ~~[[a]]~~



a
 b
 X
 -
 204 X
 ab
 X

∅: [[a]] → result in eax

100: STORE SP+4 (or +0, arg is dead here)
 101: [[b]]
 200: STORE SP+12
 201: LOAD SP+4
 202: CALL SP+8
 203: liveness vector 7
 204: STORE SP+4
 205: [[c]]
 300: STORE SP+0 (tail call)
 301: LOAD SP+4
 302: JUMP

ERROR code, action, arg,
 DECLARE q, kind, arg,
 DESTROY q
 CLONE q, acks, targetq, timeout
 PUT q, id, ctype, blob
 MISSING q, acks
 INDEX q, acks
 LIST q
 DELETE q, acks
 ? REJECT q, acks, reason

A

POST A, X
 POST B, Y
 POST B, Z

SEC
 A X
 B Y
 B Z

AES
 13 SES/1 - X

SEC

ID	ORG	TGT	BODY
1	-	AES	X
2	-	BES	Y
3	-	BES	Z

SES

1	SEC/1	AES	X
2	SEC/2	BES	Y
3	SEC/3	BES	Z

SEC		SES		AES		TR@S
1 - AES	X	1 sec/1 AES	X	13 ses/1		TR@S X

PUT SEC/1 → SES/1 AES X
 IDX ~~SES~~ ← SES {1}

nextslot(AES) = 13.

PUT SES/1 → AES/13 - X

PUT TEC/1 → TES/1 AC@8 GET(1, ∞, TR@S/15¹₂, -)

PUT AES/13 → TR@S/15 - X

over
new

new

