

par new input output rec

\emptyset
 $P \mid Q$
 $x = \lambda v. P$
 $x \Leftarrow \lambda v. P$
 $x v$
 ~~$(x) P$~~
 $x[P]$
 $\text{lift } x \Leftarrow \lambda v. P$
 $(\text{eval } v)$

stop
 par
 rep. in
 in
 out
 new
 loc.
 lift

$$\frac{\Gamma \vdash^a \lambda v. P}{\Gamma, x: \emptyset \vdash^a P} \quad x \notin \text{dn}(\Gamma)$$

$$\frac{\Gamma, x: \emptyset \vdash^a P \parallel \Delta \vdash^b Q}{\Gamma \vdash^a \lambda v. P \parallel \Delta \vdash^a Q} \quad x \notin \text{fn}(Q)$$

~~$\Gamma \vdash^a x[P]$~~

~~$\Gamma, x: \mu \vdash^a P \parallel \Delta \vdash^b (x[Q] \mid R)$~~

$$\frac{\Gamma, x: \mu \vdash^a P \parallel \Delta \vdash^b (x[Q] \mid R)}{\Gamma, x: \mu \vdash^a P \parallel \vdash^a Q \parallel \Delta \vdash^b R}$$

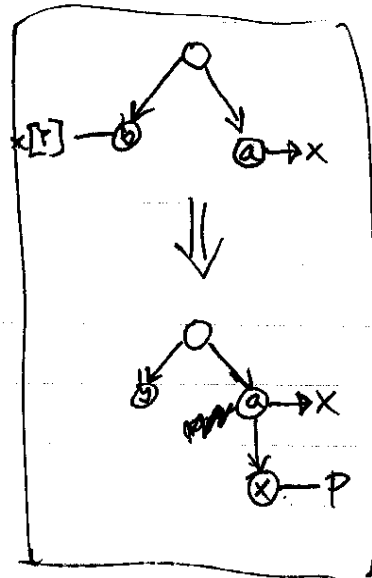
$$\Gamma \vdash^a \text{lift } x \Leftarrow \lambda v. R \parallel \Delta, x: \mu \vdash^b P \parallel \overline{\Theta \vdash^{b, x} Q}$$

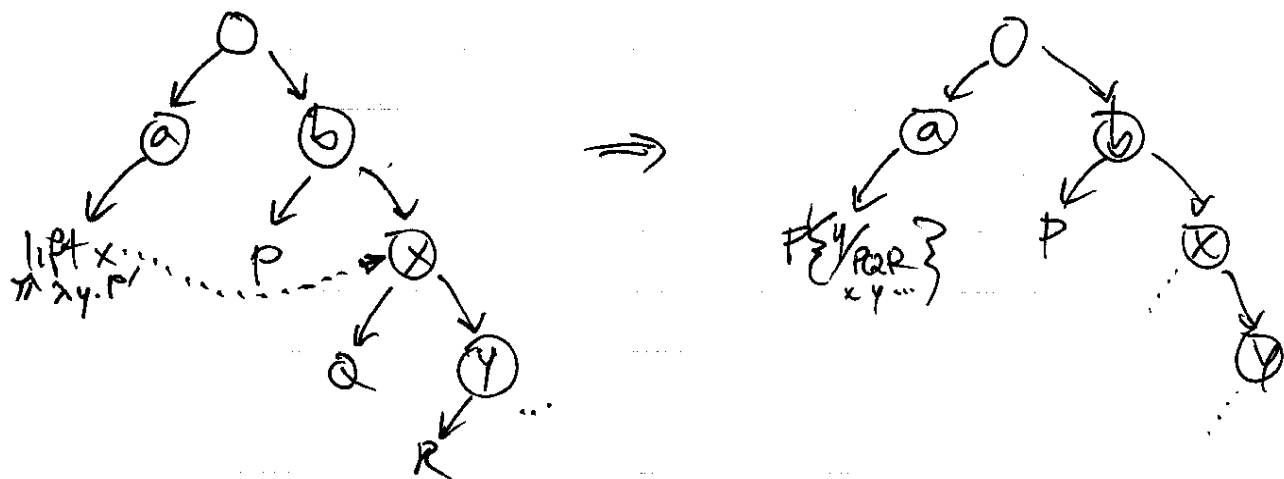
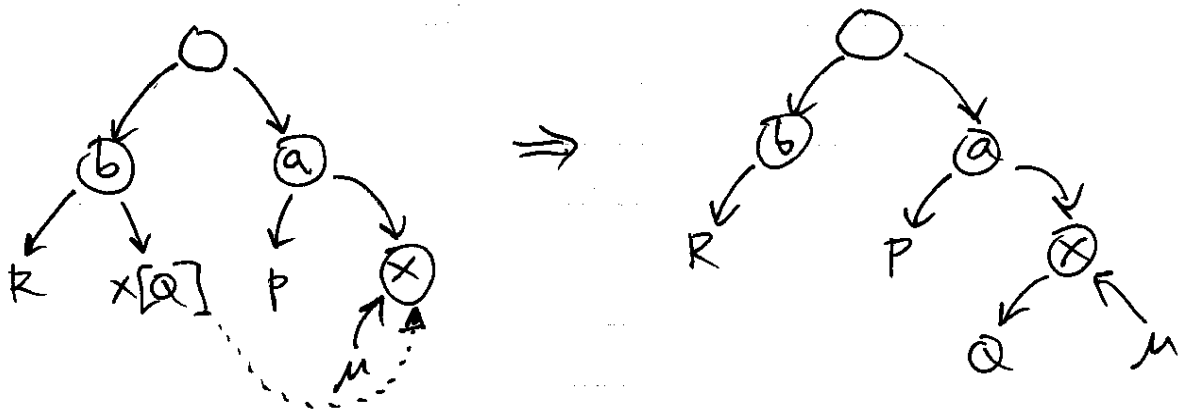
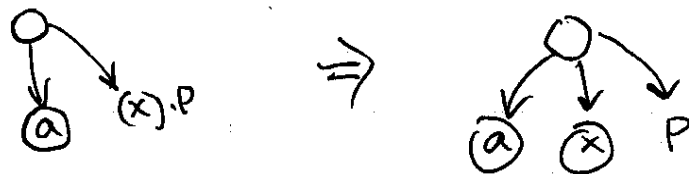
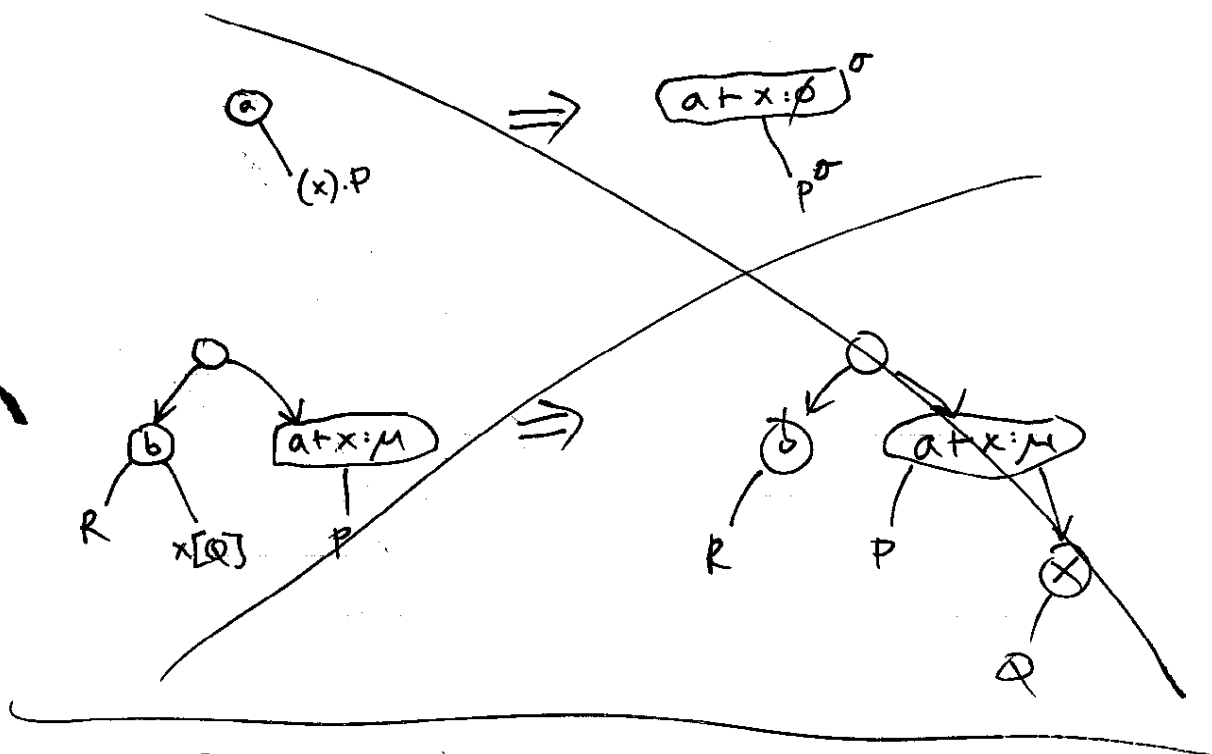
$$z: \text{LIFT}(\langle \overline{b, x}, Q \rangle) \parallel \Gamma \vdash^a (\lambda v. R) z \parallel \Delta, x: \mu \vdash^b P \parallel \overline{\Theta \vdash^{b, x} \emptyset}$$

$\text{loc} = \text{path}, \text{process}, \text{children}$
 $\text{path} = \varepsilon \mid \text{name} \cdot \text{path}$
 $\text{children} = \varepsilon \mid \text{loc}, \text{children}$
 $\text{liftresult} = \text{loc}$

-OR-

$\text{loc} = \text{name}, \text{process}, \text{parent}$
 $\text{parent} = \varepsilon \mid \text{name}$
 $\text{liftresult} = [\text{loc}]$





TREE (not DAG or graph) of locations

Since locations \equiv names,

loc = "name", process, parent, messages

... or maybe ambient-like $\langle M \rangle$ as a raw message

$$\begin{aligned} \rightarrow & \quad x \tilde{y} \Rightarrow x[\langle \tilde{y} \rangle] \\ & \quad \& x \Leftarrow \lambda \tilde{y}. P \Rightarrow x[(\tilde{y}).P] \\ & \quad \& x = \lambda \tilde{y}. P \Rightarrow x[!(\tilde{y}).P] \end{aligned}$$

... gives you input capability if you're willing to move!
— need to fix infinite nesting problem!

stop, par implicit
bipar implicit

aka $a[r[(v).a[\langle v \rangle]] \mid (v).P]$

in	$(\tilde{y}).P$	✓
rep.in.	$!(\tilde{y}).P$	✓
out	$\langle \tilde{y} \rangle$	✓
new	$\forall x. P$	✓
go	$x[P]$	✓
lift	lift $x(\tilde{y}).P$	

